Enabling Graph-Based Profiling Analysis using Hatchet

Student: Ian Lumsden

Mentors: Stephanie Brink, Michael R. Wyatt II, Todd Gamblin, and Michela Taufer

LLNL-PRES-815496

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Limits of HPC Profilers: HPCToolkit

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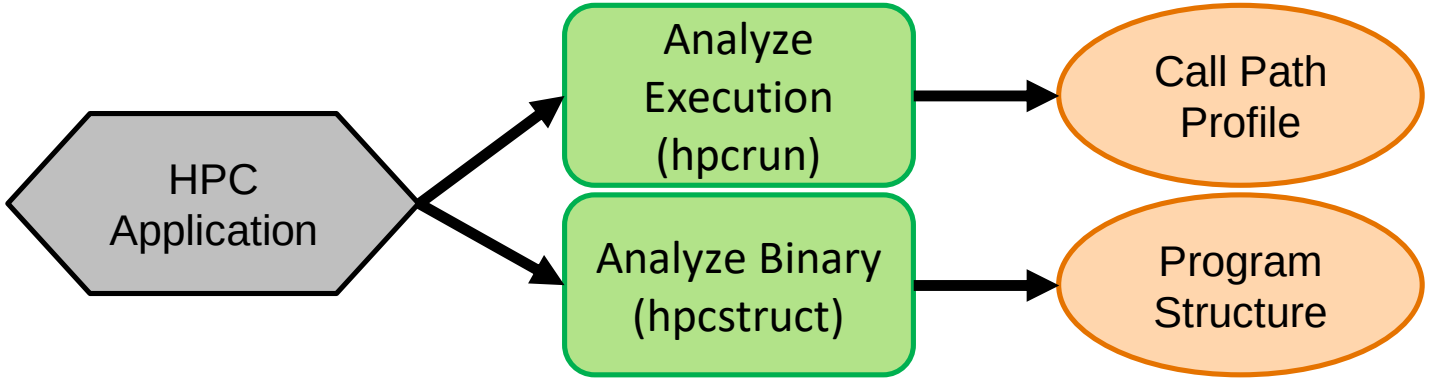

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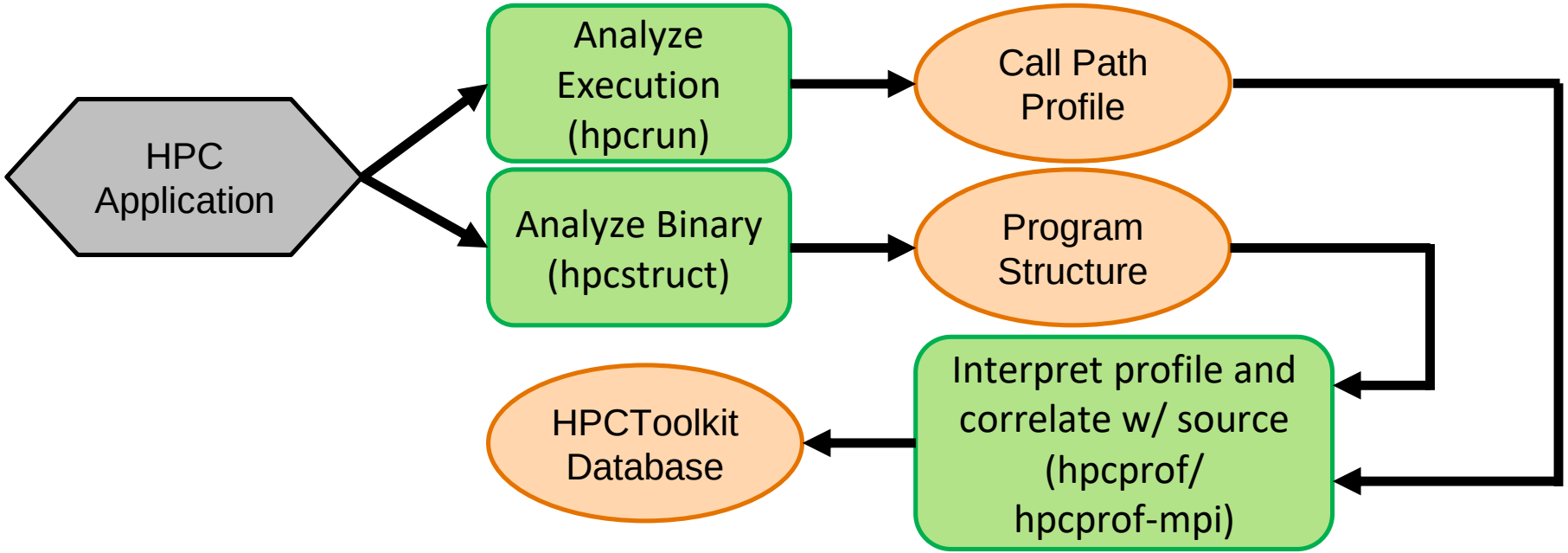

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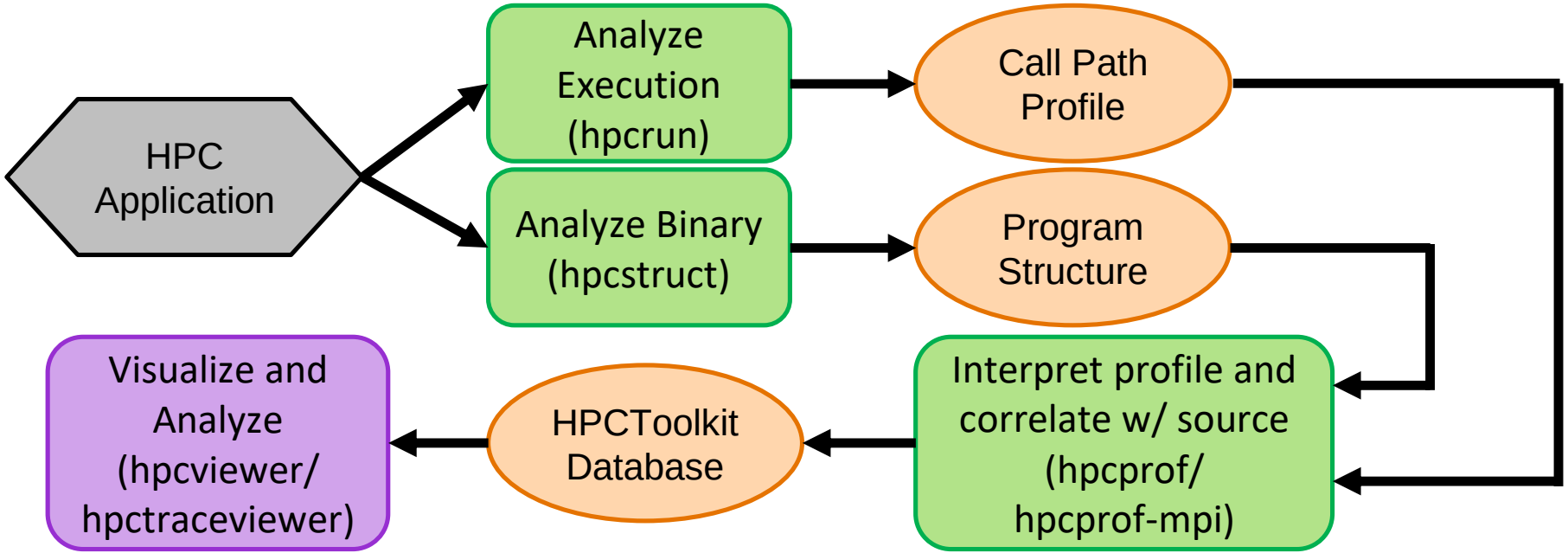

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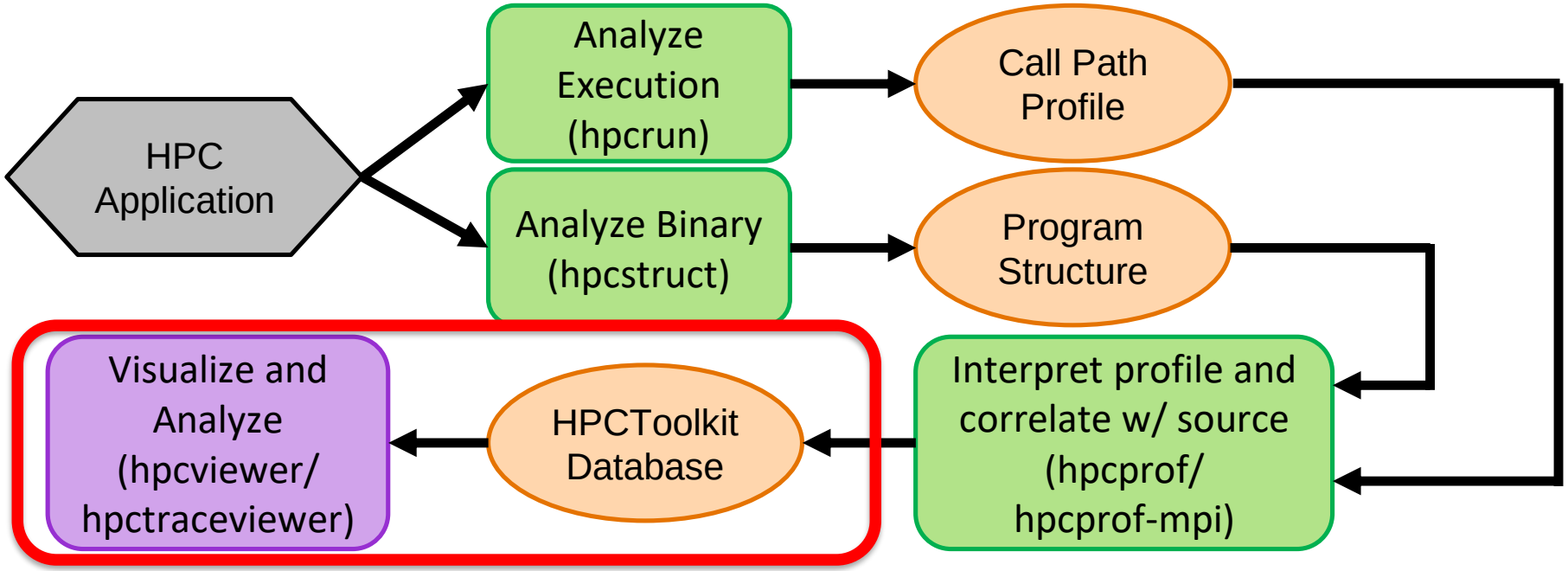

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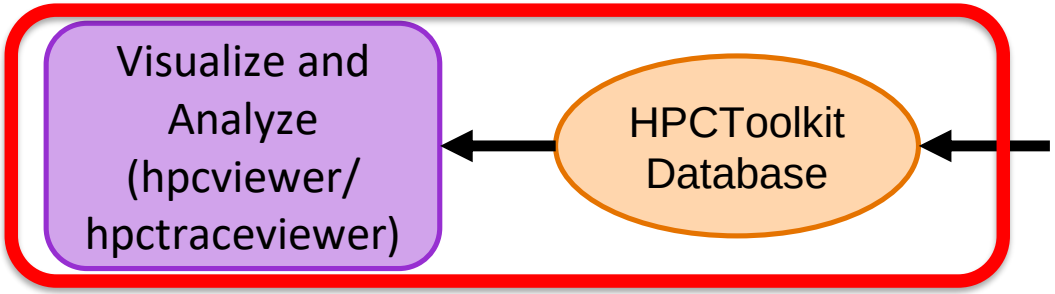

- Profiler output uses a custom, unique data format
- Visualization and analysis tools are GUI-based with limited analysis capabilities
- Both issues above limit user analysis options

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Hatchet

- Python-based library
- A new, **general** data analysis tool that can read HPC profiling data from **different profilers**
  - Store the raw performance data into a *pandas* DataFrame
  - Represent the relational *caller-callee* data with directed acyclic graph (DAG)

[1] A. Bhatele, S. Brink, and T. Gamblin. 2019. Hatchet: Pruning the Overgrowth in Parallel Profiles. In *Proc. of the International Conference of High Performance Computing, Networking, Storage, and Analysis (SC19).*

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GraphFrame

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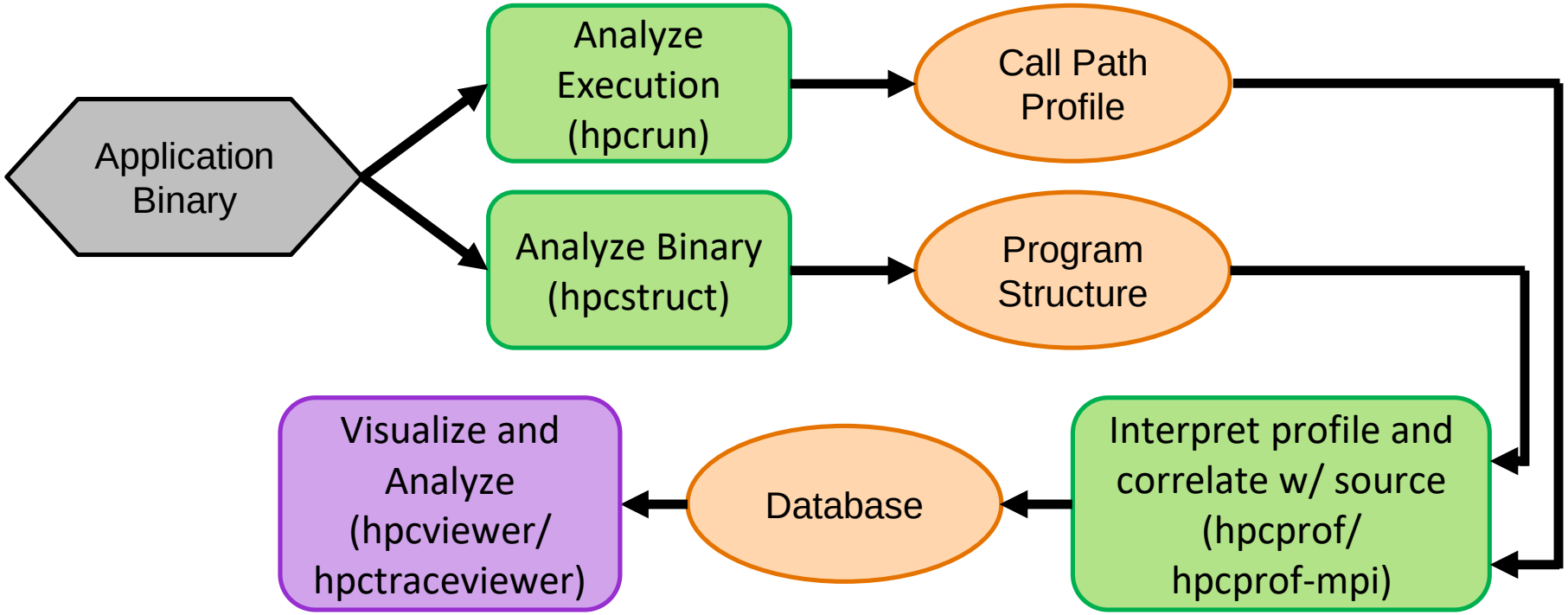
HPCToolkit with Hatchet



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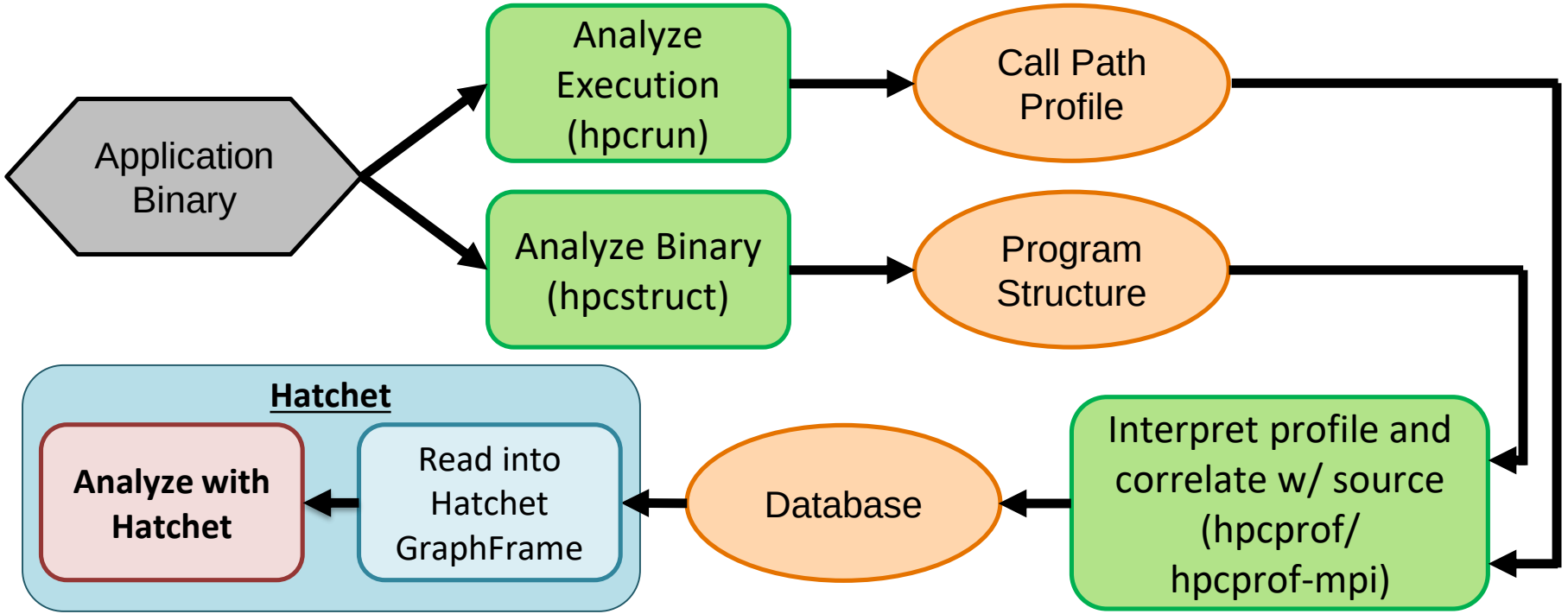

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- Hatchet restricts users to table-based analysis of the raw performance data
- Hatchet does NOT support analysis of the relational data collected by HPC profilers

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Research Goals

- Design new graph-based filtering query language to enable the use of relational data collected by profilers in analysis
- Implement this query language and integrate it into Hatchet analysis
- Use the augmented Hatchet to analyze the performance of different MPI calls in HPC benchmark applications

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    - Two parts to an *abstract graph node*:
      1. A wildcard specifying the number of real graph nodes to match to the *abstract graph node* (default is 1)
      2. A filter used to determine whether a real graph node matches the *abstract graph node* (default is an “always true” filter)

Hatchet Query Language

- Algorithm:
  1. Read and parse the user's query path
  2. Match real nodes in the graph being filtered to the abstract nodes in the query path
  3. Collect all graph paths that match the full query path
  4. Create a new graph (GraphFrame) containing only the nodes in the matched paths

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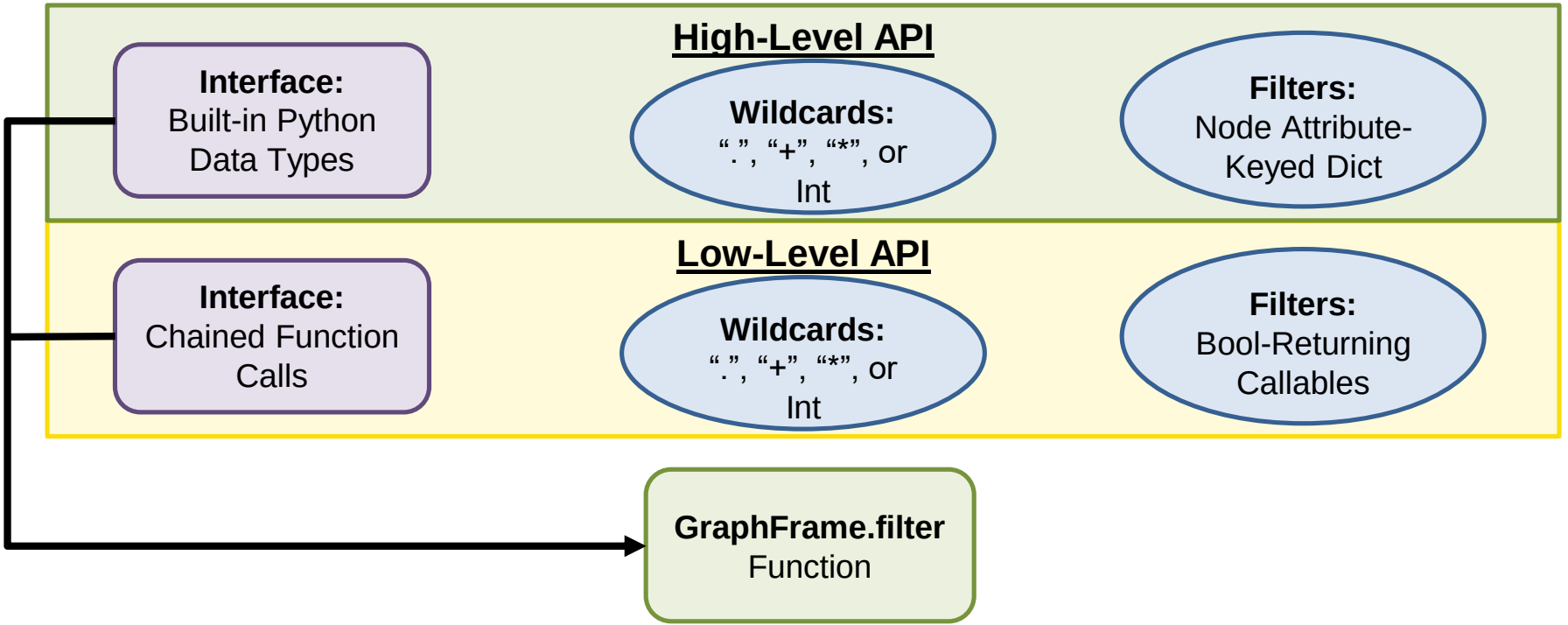
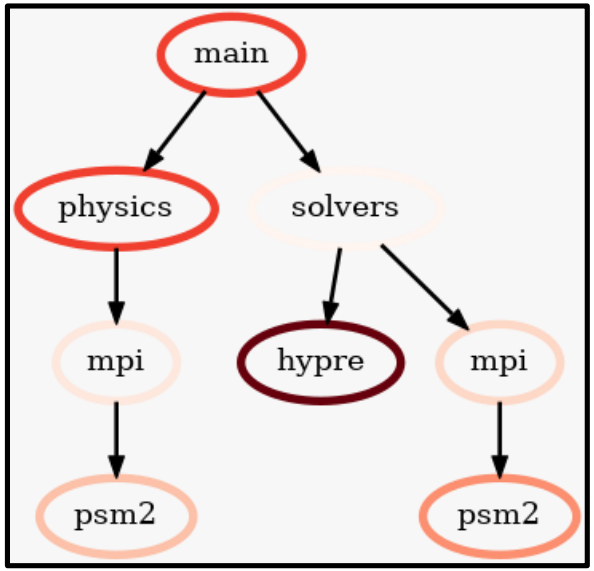
Hatchet Query Language: APIs



Hatchet Query Language: Example



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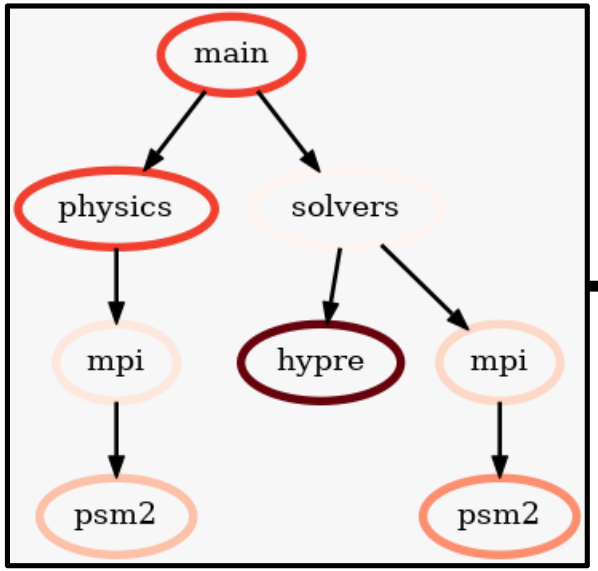


```
[
    {"name": "solvers"},
    {"time": "< 50"},
    "*"
]
```

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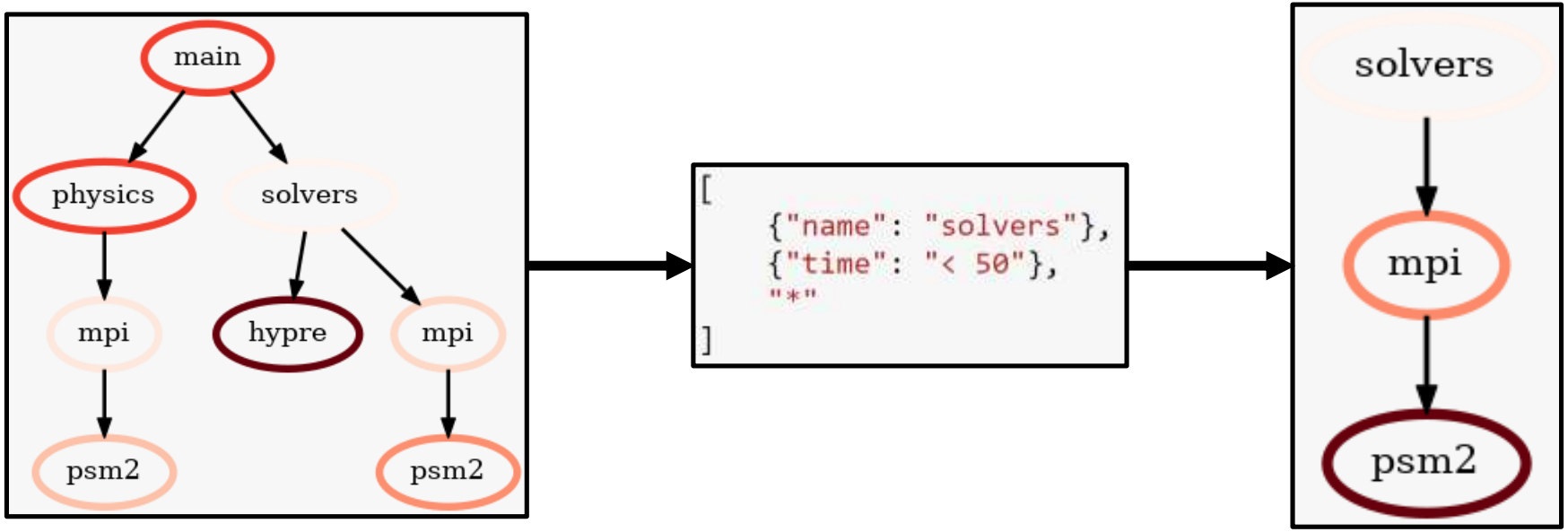

Publicly available starting with Hatchet v1.2.0

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Case Study: Collect Data

- Run each of the following benchmarks with MVAPICH (M) and Spectrum-MPI (S) using 64, 128, 256, and 512 MPI ranks, and profile the runs with HPCToolkit
  - AMG2013
  - Kripke
  - Lammps
- Load the generated profiles into Hatchet

Case Study: Collect Data

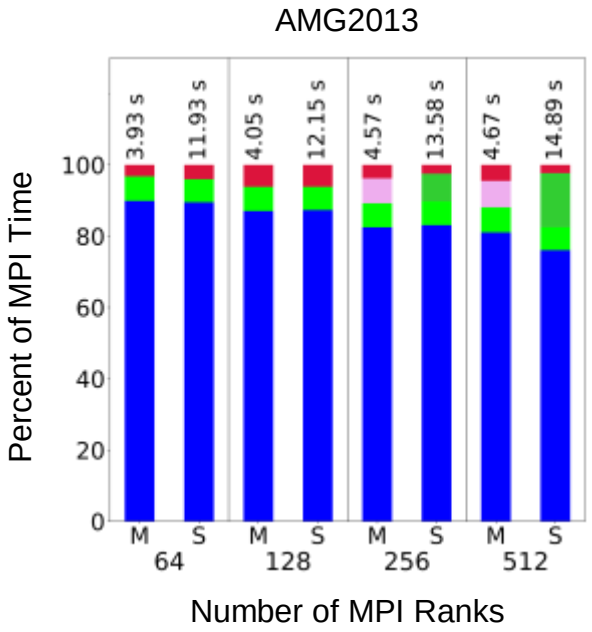
- Perform the benchmarking on LLNL's Lassen supercomputer
  - LLNL CZ (Public Collaboration) CORAL Supercomputer
  - 795 AC922 Nodes
    - 2 IBM Power9 CPUs per node (20 cores per node)
    - 256 GB Memory per node
  - Infiniband EDR interconnect

Case Study: Extract MPI Layer

- Filter the call graphs to get subgraphs rooted at standard MPI function calls
- Query Path used to extract MPI Layer:

```
[{“name”: “P?MPI_.*”}, “*”]
```

Case Study: Calculate Percent MPI Time for each MPI Function



Only MPI functions that contribute at least 5% of total MPI time are shown

MPI Function Calls			
	MPI_Finalize		MPI_Allreduce
	MPI_Allgather		MPI_Waitall
	Remaining MPI Time		

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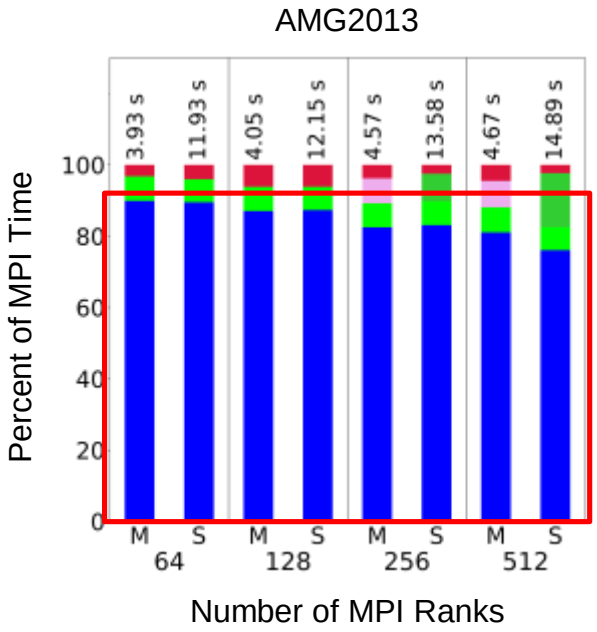

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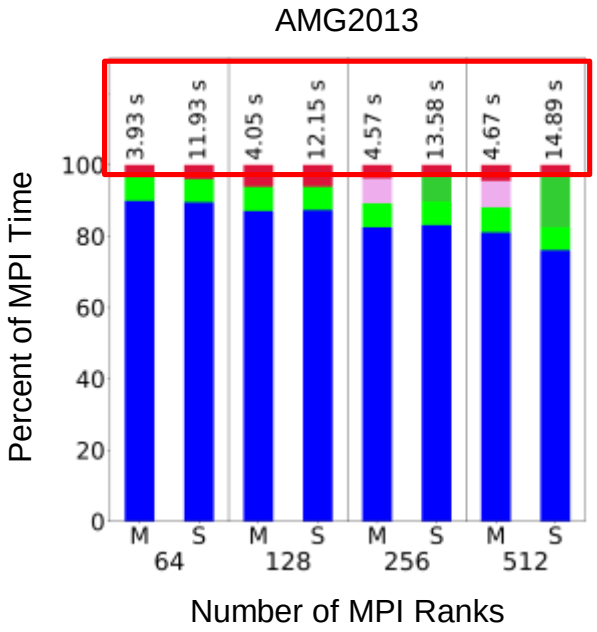

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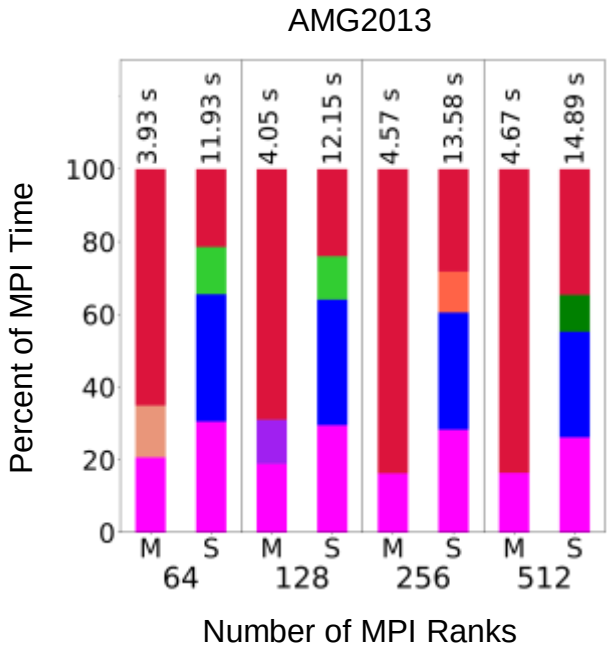

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	malloc.c:0		Geometry.h:0
	stl_vector.h:0		Remaining MPI Time

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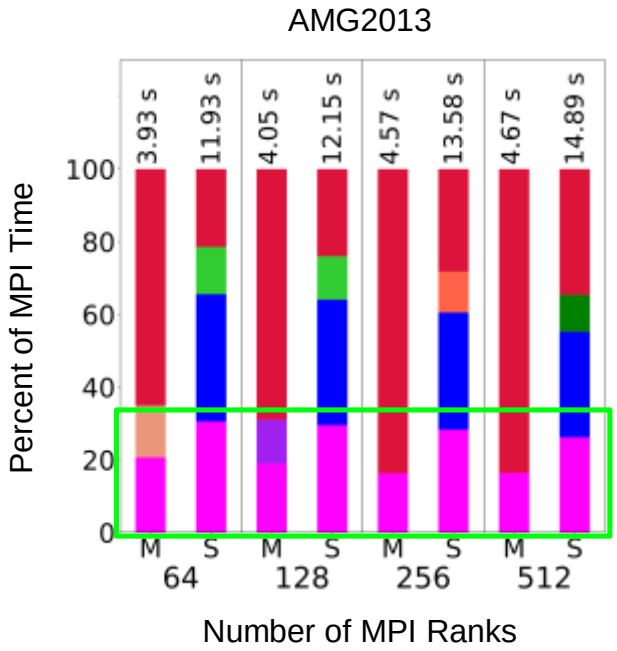

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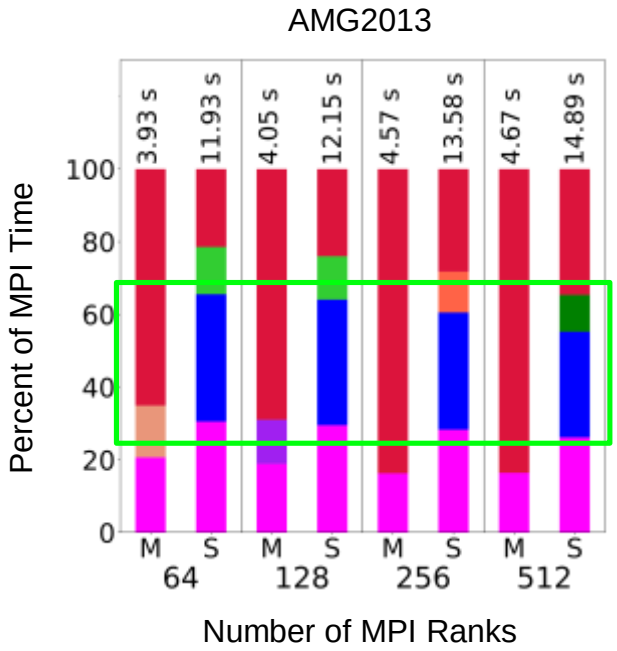

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Zoom into specific benchmarks (**AMG2013**) and examine the children of specific MPI functions (**MPI_Allgather**)

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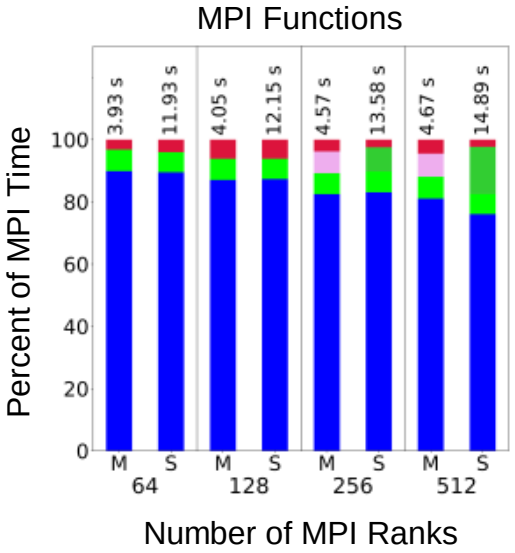

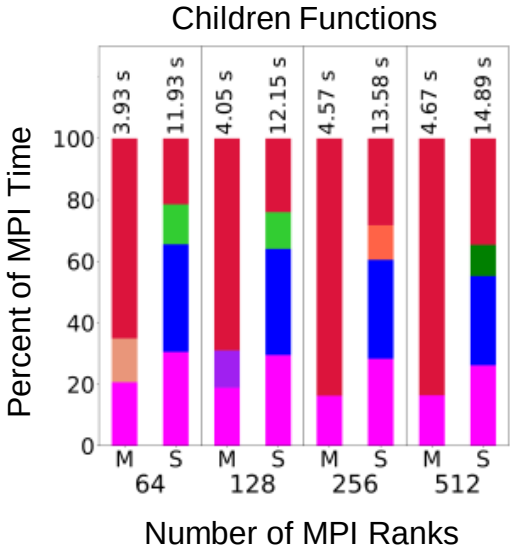

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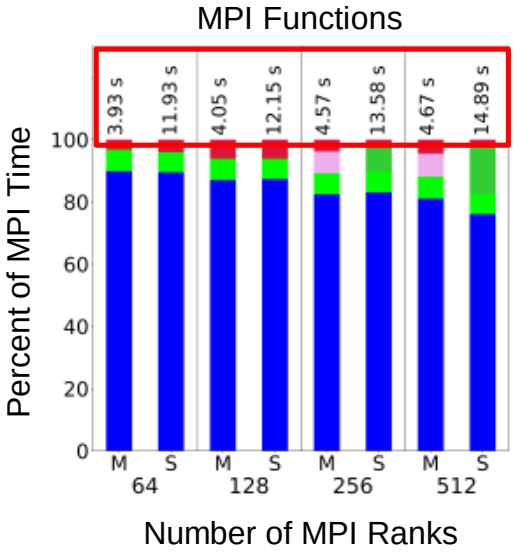

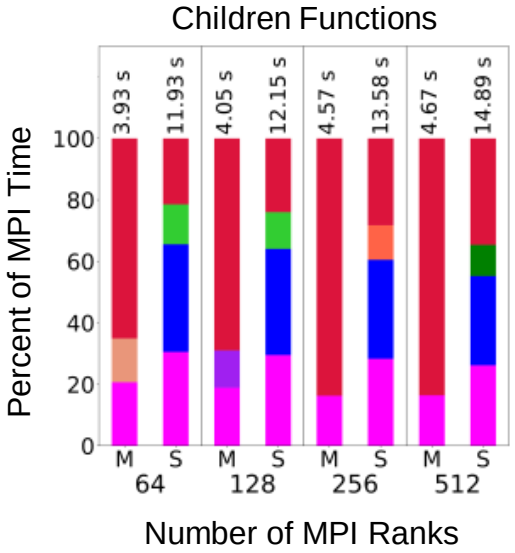

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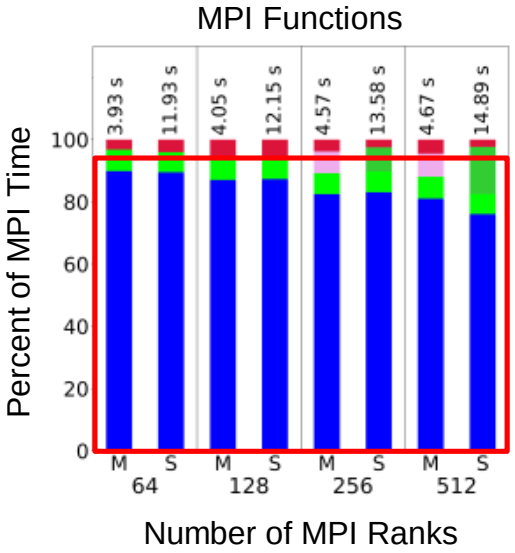

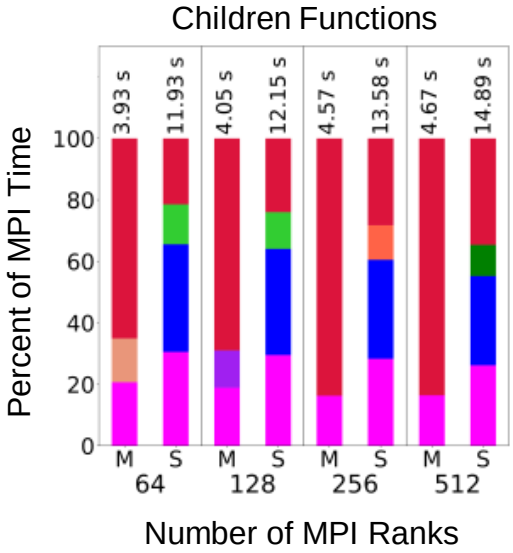

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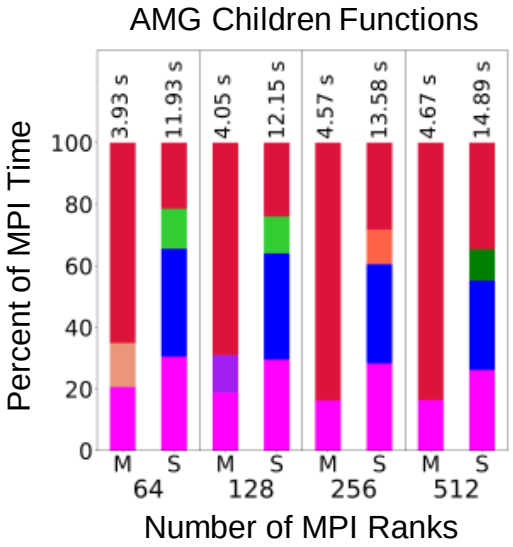

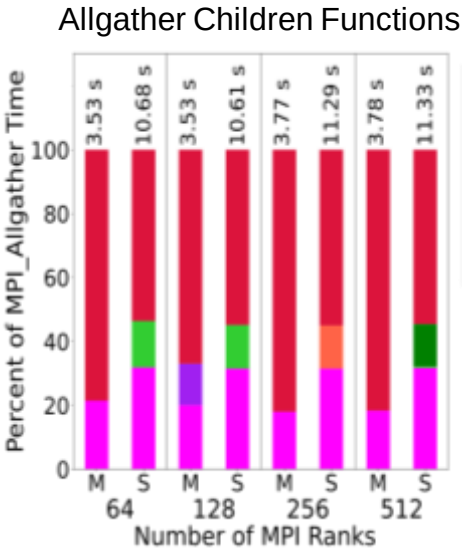

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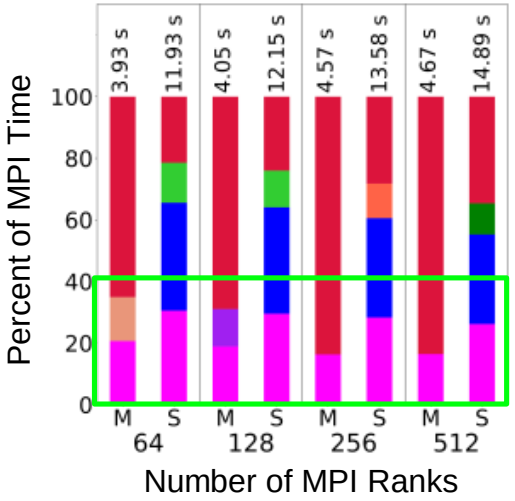
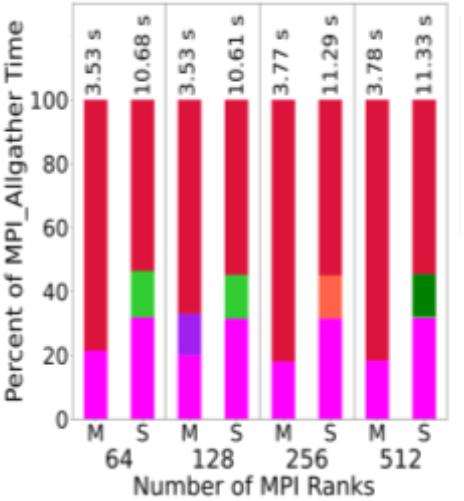
Lessons Learned: MPI

The pthread_spin_lock function is consistently a major contributor to MPI runtime (i.e., 20% or more of MPI time)

AMG Children Functions



Allgather Children Functions



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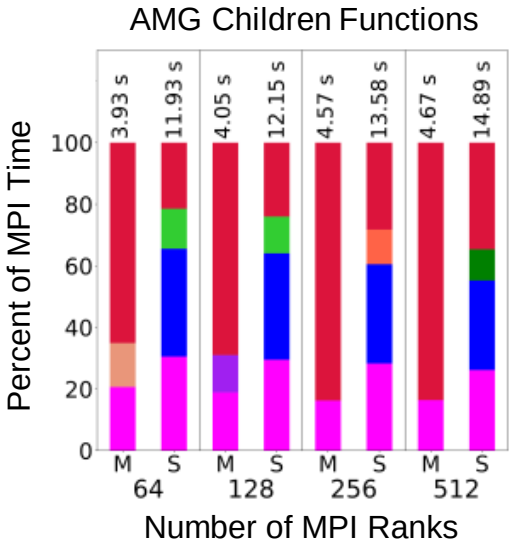
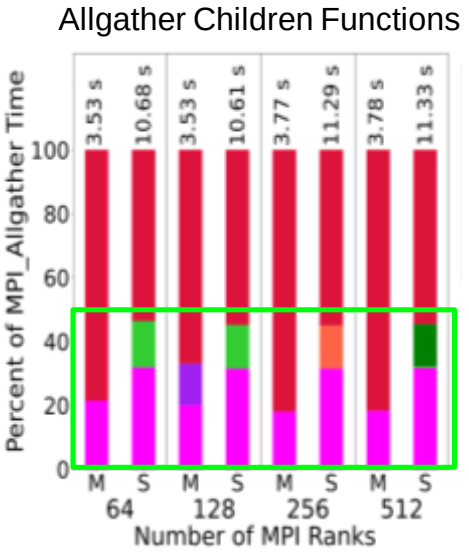
Lessons Learned: MPI

In AMG2013, the inferior performance of MPI_Allgather in Spectrum-MPI can possibly be attributed to the use of pthread_spin_lock





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Lessons Learned: Hatchet

- Using the query language and Hatchet, we are able to:
  - Extract all call paths specific to a given library
  - Determine the performance contributions of function calls used by these libraries
  - Correlate children function calls to specific important library API calls in an application
  - Use this correlation to determine children function calls that contribute the most to the performance of the targeted library API call
  - Compare the correlation of children and API calls across libraries to determine possible causes for performance differences

Future Work

- Use and expand on the techniques and tools from this work to study large scale scalability and replicability problems
  - Use information gained from studying these problems to propose mitigation strategies
  - Develop tools to help others study and mitigate these problems

Additional Resources on Hatchet

Hatchet SC19 Paper

Hatchet SC20 Paper

Hatchet GitHub

Supplemental Slides

Case Study: Benchmarks

- Kripke (LLNL) [5]
  - Mini-App Benchmark
  - Deterministic
  - Runs S_N Neutron Transport simulation

[5] AJ Kunen et al. 2015. *KRIPKE-A massively parallel transport mini-app*. Technical Report. Lawrence Livermore National Laboratory (LLNL)

Case Study: Benchmarks

- Lammps (SNL) [6]
  - Popular classical molecular dynamics simulator
  - Also contains a set of benchmarks

[6] S. Plimpton. 1995. Fast Parallel Algorithms for Short-Range Molecular Dynamics. *J. Comput. Phys.* 117, 1 (March 1995), 1-19

Case Study: Benchmarks

- AMG2013 (LLNL) [7]
  - Parallel algebraic multigrid solver for linear systems
  - Proxy App for hypre's BoomerAMG solver
  - Highly synchronous code

[7] U. Yang et al. 2017. Algebraic Multigrid Benchmark, Version 00. (8 2017)

Case Study: MPIs

- MVAPICH (v2.3.3)
  - Project led by the Network-Based Computing Laboratory at Ohio State University
  - Based on MPICH (Argonne National Lab)
  - Conforms with MPI 3.1 standard

Case Study: MPIs

- Spectrum-MPI (v10.2)
  - Created by IBM
  - Based on OpenMPI (first created by LANL, Indiana University, and the University of Tennessee)
  - Conforms with MPI 3.1 standard